

# Introduction to Transfinite Chomp

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- 1 What is Chomp?
- 2 Classical results about Chomp
- 3 Transfinite Chomp

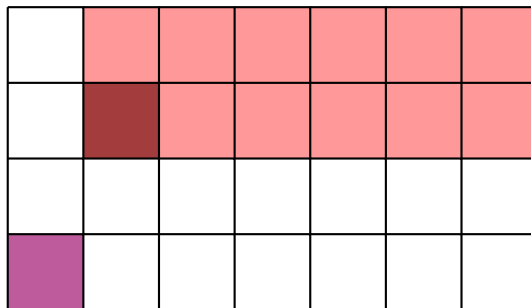
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## Before...



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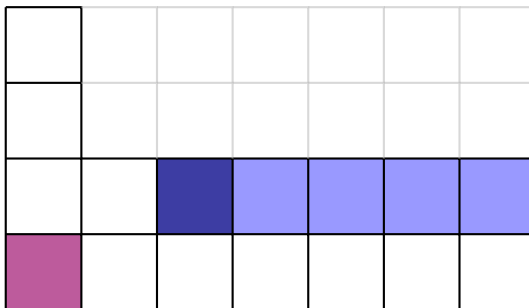
1st player's play

2nd player's play

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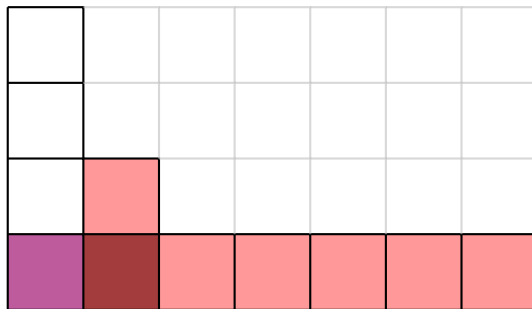
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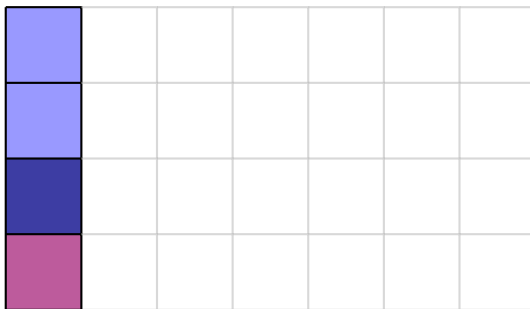
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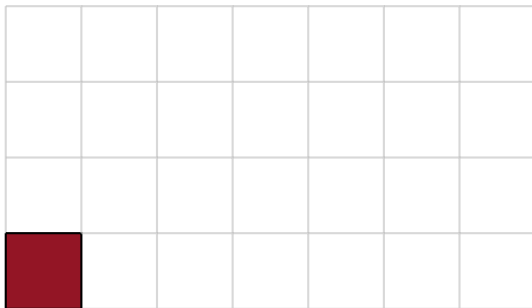
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- For which  $m, n$  is it a 1st player win?
- And what is an explicit winning strategy?

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# The strategy-stealing argument

## Theorem

$m \times n$ -chomp is a 1st player win  $\iff (m, n) \neq (1, 1)$ .

## Proof.

$\implies$ : Prove the contrapositive. It suffices to show that  $1 \times 1$ -chomp is a 2nd player win.



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This is clearly a 2nd player win.

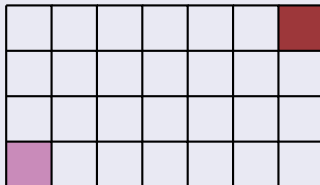
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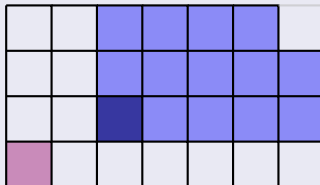
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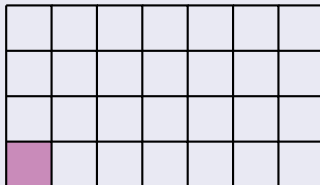
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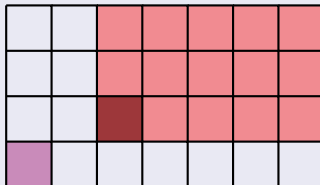
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This contradicts that it was a 2nd player win. □

# On explicit winning strategies

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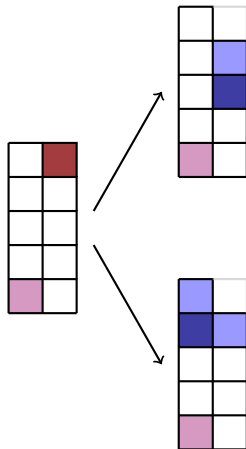
When  $(m, n) \neq (1, 1)$ , an explicit 1st player winning strategy for  $m \times n$ -chomp is generally unknown.

However, we can construct explicit winning strategies for particular cases.

## Example: $2 \times n$ -chomp

Key idea:

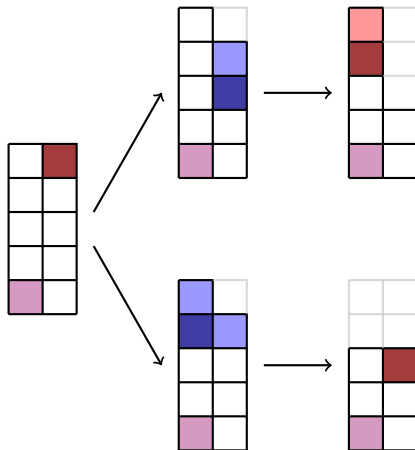
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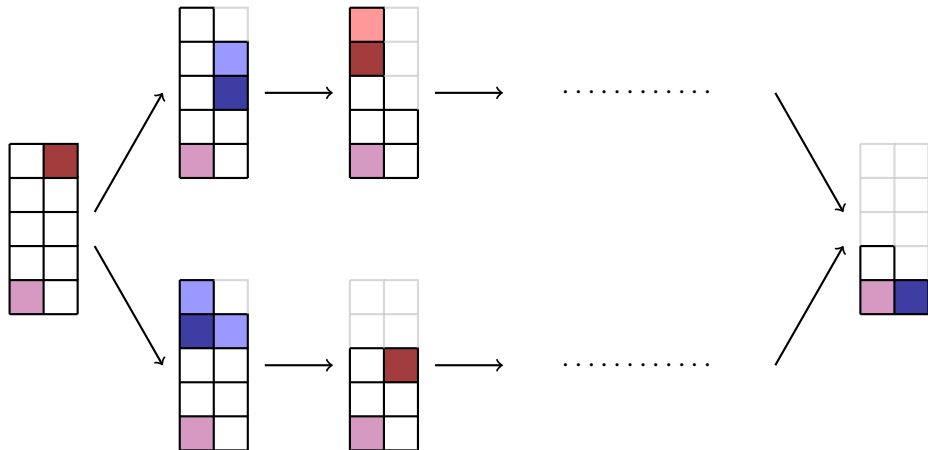
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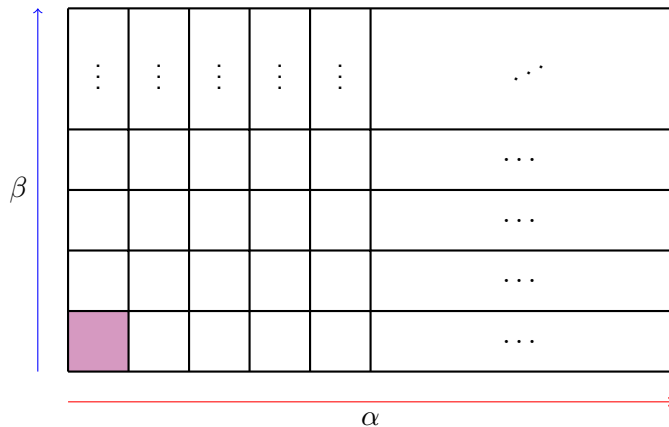


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# $\alpha \times \beta$ -chomp

For ordinals  $\alpha, \beta \geq 1$ , the  $\alpha \times \beta$ -chomp is chomp played from the following board:



# Finiteness of Transfinite Chomp

## Proposition

Transfinite chomp always ends after finitely many moves. In particular, either a 1st player or a 2nd player has a winning strategy.

Of course, constructing explicit winning strategies is very very hard.

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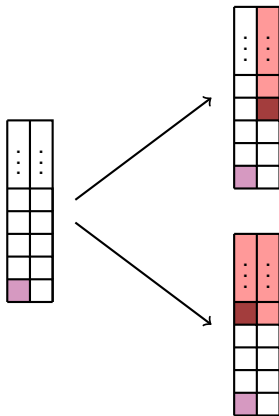
Transfinite chomp always ends after finitely many moves. In particular, either a 1st player or a 2nd player has a winning strategy.

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Also, in general, the strategy-stealing argument does not work well on  $\alpha \times \beta$ -chomp, since the upper right corner cell does not exist if  $\alpha$  or  $\beta$  limit. *Actually...*

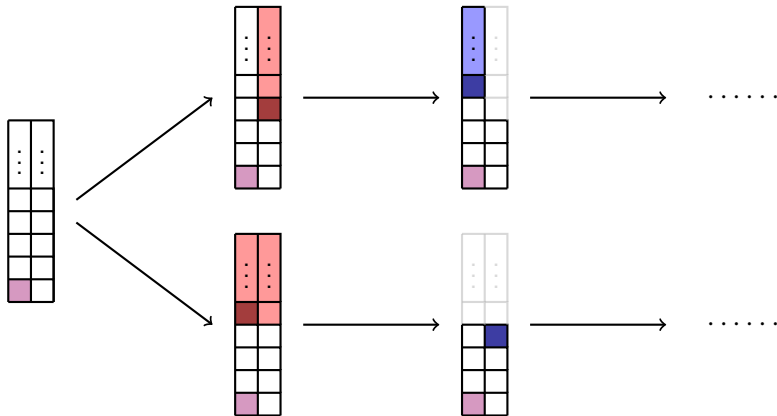
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# Class function $f$

Theorem(S. Huddleston, J. Shurman [1], Michal R. Przybylek [2] (partially), T. Goto simplified.)

For every ordinal  $\alpha \geq 1$  there exists a **unique**  $\beta \geq 1$  such that  $\alpha \times \beta$ -chomp is a 2nd player win.

e.g.,  $\omega$  is the unique ordinal  $\delta$  such that  $2 \times \delta$  is a 2nd player win.

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This defines a class function  $f$  from ordinals to ordinals by:

$$f(\alpha) = \beta \iff \alpha \times \beta \text{ is a 2nd player win.}$$

It is easy to show that  $f(\alpha) = \beta \iff f(\beta) = \alpha$ .



# Local behavior of $f$

The following are known:

$$f(1) = 1$$

$$f(2) = \omega$$

$$f(3) = \omega^\omega \quad (\text{S. Huddleston, J. Shurman[1]})$$

$$\vdots$$

$$f(\omega) = 2$$

$$\vdots$$

$$f(\omega^\omega) = 3$$

$$\vdots$$

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Let  $\text{Suc}_{\omega_1} = \{\alpha < \omega_1 \mid \alpha \text{ is a successor ordinal}\}$  and  $\text{Lim}_{\omega_1} = \{\alpha < \omega_1 \mid \alpha \text{ is a limit ordinal}\}$ . Then exactly one of  $f^{-1}(\text{Suc}_{\omega_1})$  or  $f^{-1}(\text{Lim}_{\omega_1})$  contains a club subset of  $\omega_1$ , and

$$f^{-1}(\text{Suc}_{\omega_1}) \text{ contains a club} \iff f(\omega_1) \text{ is a successor ordinal.}$$

# Reflection

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That is, the property “ $f(\alpha)$  is successor” reflects to  $\omega_1$ .

## Future works

- I would like to determine the value of  $f(\omega_1)$ .
  - At least, I would like to know whether  $f(\omega_1)$  is a successor or a limit ordinal.
  - By our second proposition, if we can compute  $f(\alpha)$  for many  $\alpha < \omega_1$ , we can answer this.

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Summary  
PLEASE HELP ME!!!

Thank you for your attention!

# References I

- [1] Scott Huddleston and Jerry Shurman.  
Transfinite chomp.  
In *More Games of No Chance, Proc. MSRI Workshop on Combinatorial Games*, pp. 183–212, 2000.
- [2] Michal R. Przybylek.  
Infinite chocolate.  
2019.