

Inner models from extended logics

(All three lectures in one file)

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Inner model from extended logics

- Lecture 1: Extended logics, inner models, examples, L -tameness.
- Lecture 2: Stationary logic, a Completeness Theorem, Club Determinacy, Applications.
- Lecture 3: Second order logic, HOD.

The tutorial is based on

- **Inner Models from Extended Logics: Part 1**, J. Kennedy, M. Magidor and J.V. Journal of Mathematical Logic (2021).
- **Inner Models from Extended Logics: Part 2**, J. Kennedy, M. Magidor and J.V. Journal of Mathematical Logic. (to appear)
- Also relevant: **Closed and unbounded classes and the Häftig quantifier model.**, Ph. Welch, J. Symb. Log. (2022).

We will learn in this first lecture:

- A new construction of a whole family of inner models.
- Why some of them are not really new.
- Why and how some of them extend known inner models.

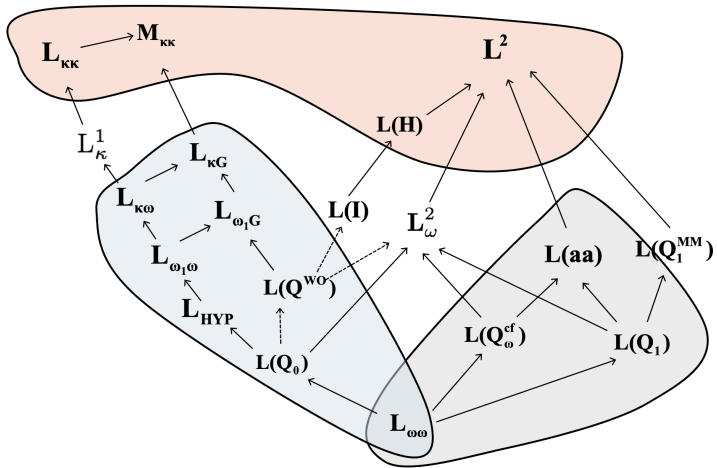


Figure: Map of extended logics

Possible desirable attributes of extended logics

- Axiomatizable.
- Downward Löwenheim-Skolem Theorem (in some form).
- Compactness Theorem (in some form).
- Can express interesting mathematical properties.
- Abstract Model Theory [1].

Common inner models

- Cumulative hierarchy V .
- Constructible sets L .
- Hereditarily ordinal definable sets HOD.
- $L[\mu]$, $L[\vec{E}]$
- $L(\mathbb{R})$
- Chang model $C_{\omega_1\omega_1}$.

Possible desirable attributes of inner models

- Forcing absolute.
- Support large cardinals.
- Arise “naturally”.
- Decide questions such as CH.
- Satisfy Axiom of Choice.

No (very) large cardinals in L .

Scott 1961 [12].

- Suppose $V = L$ and κ is (the least) MC with n.u.f. F .
- Let $N = V^\kappa$.
- Define $f \sim g \iff \{\alpha < \kappa : f(\alpha) = g(\alpha)\} \in U$.
- Define $[f] E [g] \iff \{\alpha < \kappa : f(\alpha) \in g(\alpha)\} \in U$.
- Let (M, \in) be the Mostowski collapse of (N, E) .
- Let $i(a) = [c_a]$.
- $i : V \rightarrow M$ elementary, $i(\kappa) > \kappa$.
- $M \models$ “ $i(\kappa)$ is the least MC”.
- But $M = V$, a contradiction.

Project: Inner models from extended logics

- **Replace** first order logic by one of the logics in the Map-of-Logics in order to obtain **new** inner models with desirable properties.
- The inner model $\mathcal{C}(\mathcal{L}^*)$ arises from Gödel's L by replacing first order logic $\mathcal{L}_{\omega\omega}$ by an extension \mathcal{L}^* of $\mathcal{L}_{\omega\omega}$.

The inner model L

In 1940 Gödel introduced L [3].

$$\left\{ \begin{array}{l} L_0 = \emptyset \\ L_\nu = \bigcup_{\alpha < \nu} L_\alpha \text{ for limit } \nu \\ L_{\alpha+1} = \{X \subseteq L_\alpha : X \text{ is } \mathcal{L}_{\omega\omega}\text{-definable over } L_\alpha \\ \text{i.e. } X = \{a \in L_\alpha : L_\alpha \models \varphi(a, \vec{b})\} \\ \text{for some } \varphi(x, \vec{y}) \in \mathcal{L}_{\omega\omega} \text{ and some } \vec{b} \in L_\alpha\} \\ L = \bigcup_{\alpha} L_\alpha. \end{array} \right.$$

Theorem

$L \models ZFC.$

Measuring the strengths of logics

- $\mathcal{L}^* \leq \mathcal{L}^+$ if $\mathcal{L}^* \subseteq \mathcal{L}^+$.
- $\mathcal{L}^* \leq' \mathcal{L}^+$ if $C(\mathcal{L}^*) \subseteq C(\mathcal{L}^+)$.
- A set theoretic perspective to the strength of logics.

What if we don't get anything new...

Definition

A logic \mathcal{L}^* is ***L-tame***, if $C(\mathcal{L}^*) = L$.

Can we characterize *L-tame* logics? Does *L-tameness* have model theoretic content?

Tameness of $L(Q_0)$

$Q_0 \times \varphi(x, \vec{b}) \iff$ there are infinitely many x satisfying $\varphi(x, \vec{b})$.

Theorem ([6])

$$C(\mathcal{L}(Q_0)) = L.$$

Lemma

Suppose μ is an ordinal, and $A \subseteq \mu$ such that $A \in L_\kappa$, $\kappa > \mu$. If there is a one-one $f : \omega \rightarrow A$, then there is such a function f in L_κ .

Proof.

Suppose there isn't. Since L_κ satisfies AC, there is $n < \omega$ and one-one $g : A \rightarrow n$ in L such that $g \in L_\kappa$. But such a g is a one-one $A \rightarrow n$ also in V , contradicting the existence of f in V . □

Since we can use induction on φ , the only interesting case is

$$X = \{a \in L'_\alpha : \text{There is a one-one } f : \omega \rightarrow X_a\},$$

where (by ind. hyp.)

$$X_a = \{c \in L'_\alpha : (L'_\alpha, \in) \models \psi(c, a)\} \in L_\kappa,$$

for some $\kappa > \gamma$.

Now the Lemma implies

$$X = \{a \in L'_\alpha : \text{There is a one-one } f : \omega \rightarrow X_a \text{ in } L_\kappa\},$$

Finally,

$$X = \{a \in L_\kappa : (L_\kappa, \in) \models "a \in L'_\alpha \wedge \text{There is a one-one } f : \omega \rightarrow L'_\alpha \text{ such that for all } n \psi(x, f(n)) \text{ is true when relativized to } L'_\alpha"\}.$$

Thus $X \in L_{\kappa+1}$.

More generally...

$$Q_{\alpha}x\varphi(x, \vec{b})$$



there are at least \aleph_{α} many x satisfying $\varphi(x, \vec{b})$.

- The same proof gives L -tameness of $\mathcal{L}(Q_{\alpha_0}, \dots, Q_{\alpha_n})$ for all $\alpha_0, \dots, \alpha_n$.
- $Q_{\{\alpha_1, \dots, \alpha_n\}}$ says the cardinality is one of $\{\aleph_{\alpha_1}, \dots, \aleph_{\alpha_n}\}$. This is also L -tame.

Tameness of the Magidor-Malitz logic $L(Q_0^{MM,2})$

$$Q_0^{MM,2}xy\varphi(x, y, \vec{b}) \iff$$

there is an **infinite** set X such that $\forall x, y \in X \varphi(x, y, \vec{b})$.

Theorem ([6])

$$C(Q_0^{MM,2}) = L.$$

Lemma

Suppose μ is an ordinal, and $A \in L_\kappa$, $\kappa > \mu$, such that $A \subseteq [\mu]^2$. If there is an infinite B such that $[B]^2 \subseteq A$, then there is such a set B in L_κ .

Proof.

Blackboard! □

Tameness of the Magidor-Malitz quantifier, assuming 0^\sharp

$Q_1^{MM}xy\varphi(x, y, \vec{b}) \iff$
 there is an **uncountable** set X such that $\forall x, y \in X \varphi(x, y, \vec{b})$.

This is **not** L -tame in general, but:

Theorem ([6])

If 0^\sharp exists, then $C(Q_1^{MM}) = L$.

The story of 0^\sharp

“ 0^\sharp exists” means and implies that there is a club class I of ordinals such that

1. $L \models \varphi(\gamma_1, \dots, \gamma_n) \leftrightarrow \varphi(\gamma'_1, \dots, \gamma'_n)$ whenever $\gamma_1 < \dots < \gamma_n, \gamma'_1 < \dots < \gamma'_n$ are in I and $\varphi(x_1, \dots, x_n)$ is a first order formula of set theory.
2. $L_\gamma \prec L$ whenever $\gamma \in I$.
3. If $\gamma \in \text{Lim}(I)$, then $\{X \subseteq \gamma : \exists \delta((I \setminus \delta) \cap \gamma \subseteq X)\}$ is an L -ultrafilter.
4. **Rowbottom Property:** Suppose $\gamma \in \text{Lim}(I)$. Suppose $C \subseteq [\gamma]^2$, where $C \in L$. Then there is $B \in \mathcal{U}_\gamma$ such that $[B]^2 \subseteq C$ or $[B]^2 \cap C = \emptyset$.

Lemma

Suppose 0^\sharp exists, μ is an ordinal, and $A \in L$ such that $A \subseteq [\mu]^2$.
 If there is an uncountable B such that $[B]^2 \subseteq A$, then there is such a set B in L .

How does the Lemma helps us to prove the theorem

- We will use induction on α to prove that $L'_\alpha \subseteq L$.
- We suppose $L'_\alpha \subseteq L$, and hence $L'_\alpha \in L_\gamma$ for some canonical indiscernible γ .
- We show that $L'_{\alpha+1} \subseteq L_{\gamma+1}$.
- Suppose $X \in L'_{\alpha+1}$.
- Then X is of the form

$$X = \{a \in L'_\alpha : (L'_\alpha, \in) \models \varphi(a, \vec{b})\},$$

where $\varphi(x, \vec{y}) \in \mathcal{L}(Q_1^{MM})$ and $\vec{b} \in L'_\alpha$.

- For simplicity we suppress the mention of \vec{b} .

The Lemma implies

$$X = \{a \in L'_\alpha : \text{There is an uncountable set } Y \subseteq L'_\alpha, Y \in L \text{ such that } [Y]^2 \subseteq X_a\},$$

Since $L_\gamma \prec L$, we have

$$X = \{a \in L'_\alpha : \text{There is an uncountable set } Y \subseteq L'_\alpha, Y \in L_\gamma \text{ such that } [Y]^2 \subseteq X_a\},$$

Finally,

$$X = \{a \in L_\gamma : (L_\gamma, \in) \models "a \in L'_\alpha \wedge \text{There is an uncountable } Y \subseteq L'_\alpha \text{ such that for all } x, y \in Y \psi(x, y, a) \text{ is true when relativized to } L'_\alpha" \} \in L_{\gamma+1}.$$

Let

$$C = \{ \{\alpha, \alpha'\} \in [\gamma]^2 : \{\tau(\alpha), \tau(\alpha')\} \in A \}. \quad (1)$$

Since $A \in L$, also $C \in L$. By the Rowbottom Property there is $B_0 \in \mathcal{U}_\gamma$ such that

$$[B_0]^2 \subseteq C \text{ or } [B_0]^2 \cap C = \emptyset. \quad (2)$$

Claim: $[B_0]^2 \subseteq C$.

To prove this suppose $[B_0]^2 \cap C = \emptyset$. Since $B_0 \in \mathcal{U}_\gamma$, there is $\xi < \gamma$ such that $(I \setminus \xi) \cap \gamma \subseteq B_0$. We can now find $i, j < \omega_1$ such that

$$\xi < \alpha^i < \gamma, \xi < \alpha^j < \gamma.$$

Then since by the choice of B ,

$$\tau(\alpha^i), \tau(\alpha^j) \in B,$$

and $[B]^2 \subseteq A$, we have

$$\{\tau(\alpha^i), \tau(\alpha^j)\} \in A.$$

Hence

$$\{\alpha^i, \alpha^j\} \in C \tag{3}$$

contrary to the assumption $[B_0]^2 \cap C = \emptyset$. We have proved the claim.

A case of non-tameness

Shelah [13] introduced:

$$Q_\omega^{\text{cof}} xy\varphi(x, y, \vec{b}) \iff$$

$\varphi(x, y, \vec{b})$ defines a linear order of countable cofinality.

Define $C^* = C(\mathcal{L}(Q_\omega^{\text{cof}}))$.

Not L -tame:

Theorem ([6])

If 0^\sharp exists, then $0^\sharp \in C^*$, so $C^* \neq L$.

- Let I be the canonical set of indiscernibles obtained from 0^\sharp .
- CLAIM: ordinals ξ which are regular cardinals in L and have cofinality $> \omega$ in V are in I .
- Suppose $\xi \notin I$. Note that $\xi > \min(I)$.
- Let δ be the largest element of $I \cap \xi$.
- Let $\lambda_1 < \lambda_2 < \dots$ be an infinite sequence of elements of I above ξ .

- Since ξ is regular in L , $\eta_n < \xi$.
- Since ξ has cofinality $> \omega$ in V , $\eta = \sup_n \eta_n < \xi$.
- But we have now proved that every $\alpha < \xi$ is below η , a contradiction.
- So we may conclude that necessarily $\xi \in I$.

Let

$$X = \{\xi \in L'_{N_\omega} : (L'_{N_\omega}, \in) \models \text{"}\xi \text{ is regular in } L" \wedge \neg Q_\omega^{cf} xy (x \in y \wedge y \in \xi)\}$$

Now X is an infinite subset of I and $X \in C(Q_\omega^{cof})$. Hence $0^\sharp \in C(Q_\omega^{cof})$:

$$0^\sharp = \{\ulcorner \varphi(x_1, \dots, x_n) \urcorner : (L_{N_\omega}, \in) \models \varphi(\gamma_1, \dots, \gamma_n)$$

for some $\gamma_1 < \dots < \gamma_n$ in $X\}$.

More about C^* later.

A summary of this lecture

- Abstract logics \mathcal{L}^* give rise to inner models $C(\mathcal{L}^*)$.
- Some logics are provably L -tame.
- Some logics (Magidor-Malitz logic) are L -tame assuming 0^\sharp .
- Some logics (cofinality logic) are not L -tame assuming 0^\sharp .

Next lecture

- Stationary logic and its inner model $C(aa)$.
- A Completeness Theorem using iterated generic ultrapowers.
- Club Determinacy from a proper class of Woodin cardinals.
- Applications: forcing absoluteness, large cardinals, CH

This lecture

- Stationary logic and its inner model $C(\text{aa})$.
- A Completeness Theorem using iterated generic ultrapowers. (Joint work with B. Velickovic)
- Club Determinacy from a proper class of Woodin cardinals.
- Applications: forcing absoluteness, large cardinals, CH

Stationary logic $L(\text{aa})$

Suppose we have a model \mathcal{M} and a formula $\varphi(s, \vec{b})$.

$$\begin{array}{c|ccc}
 I & x_0 & x_2 & \dots \\
 \hline
 II & & x_1 & x_3 \dots
 \end{array}$$

Player II wins if the set $\{x_0, x_1, x_2, \dots\}$ satisfies $\varphi(s, \vec{b})$ in \mathcal{M} .

We write this $\mathcal{M} \models \text{aa}s\varphi(s, \vec{b})$.

“aa” is short for “almost all”.

Stationary logic $L(\mathbf{aa})$

$$\begin{aligned}
 \mathbf{aa}s\varphi(s, \mathbf{a}) &\iff \forall x_0 \exists x_1 \forall x_2 \exists x_3 \dots \varphi(\{x_0, x_1, x_2, \dots\}, \mathbf{a}) \\
 &\iff \{A \in \mathcal{P}_{\omega_1}(M) : (\mathcal{M}, A) \models \varphi(A, \mathbf{a})\} \\
 &\quad \text{contains a club of countable subsets } M. \\
 \\
 Q_1 x \varphi(x, \mathbf{a}) &\iff |\{b \in M : \mathcal{M} \models \varphi(b, \mathbf{a})\}| \geq \aleph_1 \\
 &\iff \neg \mathbf{aa}s \forall y (\varphi(y) \rightarrow s(y)). \\
 \\
 Q_\omega^{\text{cof}} xy \varphi(x, y, \mathbf{a}) &\iff \mathbf{aa}s \forall x (\exists y \varphi(x, y) \rightarrow \exists y (\varphi(x, y) \wedge s(y)))
 \end{aligned}$$

A completeness theorem¹ for $\mathcal{L}(\text{aa})$

- Consider
 - (1) φ has a model and
 - (2) there is a model (of set theory) for “ φ has a model”.
- We prove (1) and (2) are equivalent.
- The easy direction is (1) \implies (2). Follows from Reflection.
- To prove the other direction we start with a countable model M_0 of “ φ has a model \mathcal{A} ”.
- We construct an “ $\mathcal{L}(\text{aa})$ -absolute” (to be explained) elementary extension M_{ω_1} of M_0 .
- Then M_{ω_1} satisfies “ φ has a model $j_{\omega_1}(\mathcal{A})$ ”.
- Since M_{ω_1} is “ $\mathcal{L}(\text{aa})$ -absolute”, φ really has a model, and we have (1).

¹[2], [13]

A completeness theorem for $\mathcal{L}(aa)$

- We build an elementary chain of length ω_1^V .
- $M_0 \xrightarrow{j_{01}} M_1 \xrightarrow{j_{12}} M_2 \rightarrow \cdots \rightarrow M_\alpha \xrightarrow{j_{\alpha\alpha+1}} M_{\alpha+1} \cdots \rightarrow M_{\omega_1}$
- Notation: $j_{0\alpha} = j_\alpha$.

A completeness theorem for $\mathcal{L}(\text{aa})$

- $j_1(\alpha) = \text{“}\alpha\text{”}$ i.e. has exactly α predecessors in M_1 , when $\alpha < \omega_1$.
- $\sup_{\alpha < \omega_1^M} (j_1(\alpha)) = [id]$
- $j_1(\omega_1^M) > [id]$
- A new element $[id]$ is put to the ω_1 of the model M_1 .
- Consequence: $\omega_1^{M_{\omega_1}}$ is \aleph_1 -like.

Completeness Theorem for $L(aa)$

- Some hassle arises from the fact that $L(aa)$ has second order variables in addition to first order variables.
- If the domain of the model is ω_1 (or just \aleph_1 -like with a copy of ω_1 inside, call it E), we can aa-quantify over **countable ordinals** (or initial segments determined by elements of E in the case of an \aleph_1 -like ordering with a copy E of ω_1 inside) rather than countable subsets.
- This is because if C is a club of countable subsets of ω_1 , then the set $D \subseteq C$ consisting of countable ordinals (or initial segments determined by elements of E in the case of an \aleph_1 -like ordering with a copy E of ω_1 inside) that are (as sets) in C is also a club as a set of countable ordinals.

Pulling $j_{\omega_1}(\mathcal{A})$ from M_{ω_1} to V

- $M_0 \models " \mathcal{A} \models \varphi "$, and therefore $M_{\omega_1} \models " j_{\omega_1}(\mathcal{A}) \models \varphi "$
- We define a model \mathcal{A}^* in V .
- The domain of \mathcal{A}^* is $A^* = \{ a \in M_{\omega_1} : M_{\omega_1} \models a \in \omega_1^{M_{\omega_1}} \}$
- $R^{\mathcal{A}^*} = \{ (a, b) \in A^* \times A^* : M_{\omega_1} \models " j_{\omega_1}(\mathcal{A}) \models R(a, b) " \}$

The goal

1. Starting point is $\mathcal{A} \models \varphi$, where $\varphi \in L(\mathbf{aa})$.
2. We have reached: $M_{\omega_1} \models "j_{\omega_1}(\mathcal{A}) \models \varphi"$.
3. We need to prove by induction on subformulas ψ of φ :

$$(\star) \quad \mathcal{A}^* \models \psi \iff M_{\omega_1} \models "j_{\omega_1}(\mathcal{A}) \models \psi".$$

4. Then, letting ψ be φ , $\mathcal{A}^* \models \varphi$ will follow from (2) and (3) and we are done.
5. As usual, we need to add **parameters** to (\star) as they arise from quantifiers.
6. To deal with parameters we adopt a lot constant symbols.

Some preliminaries

- Let K be a set of \aleph_1 new constant symbols. Assume $K \in M_0$.
- Let $\{S_{\varphi(s, \vec{x}), \vec{c}} : \varphi(s, \vec{x}) \in L(\text{aa}), \vec{c} \in K^{<\omega}\}$ be a splitting in V of ω_1 into ω_1 disjoint stationary sets.
- When we construct the models M_α we make sure that every element of M_α which is in $\omega_1^{M_\alpha}$ is the value of a constant symbol from K .
- That is, we keep expanding the models M_α so that every element of M_α which is in $\omega_1^{M_\alpha}$ is the value of a constant symbol from K .
- Since always the domain of $j_\alpha(\mathcal{A})$ is $\omega_1^{M_\alpha}$, and we consider aa-truth in $j_\alpha(\mathcal{A})$ only, we may drop the parameters, because they are represented by the constant symbols.

A completeness theorem for $\mathcal{L}(\text{aa})$

- Suppose M_{ω_1} satisfies “ $j_{\omega_1}(\mathcal{A}) \models \text{aa } s\psi(s)$ ”, and, for simplicity, there are no parameters.
- Then M_0 satisfies “ $\mathcal{A} \models \text{aa } s\psi(s)$ ”.
- Hence $\{\alpha < \omega_1^{M_0} : M_0 \models “\mathcal{A} \models \psi'(\alpha)”\}$ is a club and therefore is in G_0 , **whichever** way G_0 is chosen. Here $\psi'(\alpha)$ is obtained from $\psi(s)$ by changing everywhere “ $s(t)$ ” to $t < \alpha$.
- Hence $M_1 \models “\mathcal{A} \models \psi'([id_0])”$.
- Similarly, $M_{\alpha+1} \models “j_{\alpha+1}(\mathcal{A}) \models \psi'([id_\alpha])”$ for all α .
- In the end, $M_{\omega_1} \models “j_{\omega_1}(\mathcal{A}) \models \psi'(j_{\alpha\omega_1}[id_\alpha])”$ for all α .
- By Ind. Hyp, $\mathcal{A}^* \models \psi'(j_{\alpha\omega_1}[id_\alpha])$ for all α .
- Hence $\mathcal{A}^* \models \text{aa } s\psi(s)$.

A completeness theorem for $\mathcal{L}(\text{aa})$

- Conversely, suppose M_{ω_1} satisfies “ $j_{\omega_1}(\mathcal{A}) \not\models \text{aa } s\psi(s)$ ”.
- Thus M_{ω_1} satisfies “ $j_{\omega_1}(\mathcal{A}) \models \text{stat } s\neg\psi(s)$ ”.
- Then M_0 satisfies $\mathcal{A} \models \text{stat } s\neg\psi(s)$.
- Hence $S = \{\alpha < \omega_1^{M_0} : M_0 \models “\mathcal{A} \models \neg\psi'(\alpha)”\}$ is stationary and therefore we can choose G_0 so that S is in G_0 .
- Hence $M_1 \models “\mathcal{A} \models \neg\psi'([id_0])”$.
- Similarly, $M_{\alpha+1} \models “j_{\alpha+1}(\mathcal{A}) \models \neg\psi'([id_\alpha])”$ for all α . But we choose G_α so that (the corresponding set) S is in G_α **only if** $\alpha \in S_{\psi(s), \vec{c}}$, where \vec{c} are the constant occurring in $\psi(s)$.
- In the end, $M_{\omega_1} \models “j_{\omega_1}(\mathcal{A}) \models \neg\psi'(j_{\alpha\omega_1}[id_\alpha])”$ for all $\alpha \in S_{\psi'(s), \vec{c}}$.
- By Ind. Hyp, $\mathcal{A}^* \models \neg\psi'(j_{\alpha\omega_1}[id_\alpha])$ for all $\alpha \in S_{\psi'(s), \vec{c}}$.
- Hence $\mathcal{A}^* \models \text{stat } s\neg\psi(s)$.
- Hence $\mathcal{A}^* \not\models \text{aa } s\psi(s)$.

The definition of $C(\text{aa}) = C(\mathcal{L}(\text{aa}))$

Because $\mathcal{L}(\text{aa})$ has second order variables, our definition of $C(\mathcal{L}(\text{aa}))$ does not guarantee that $C(\mathcal{L}(\text{aa}))$ satisfies Axiom of Choice. It is an **open problem** whether it does. Therefore we modify the construction. We do not know whether it is a proper modification. Another **open problem!**

Jensen's J -hierarchy:

Suppose T is a class.

$$\begin{cases} J_0^T & = \emptyset \\ J_{\alpha+1}^T & = \text{rud}_T(J_\alpha^T \cup \{J_\alpha^T\}) \\ J_\nu^T & = \bigcup_{\alpha < \nu} J_\alpha^T, \text{ for } \nu = \bigcup \nu. \end{cases}$$

Here rud_T includes the operation $x \mapsto x \cap T$.

The definition of $C(\mathbf{aa}) = C(\mathcal{L}(\mathbf{aa}))$

We define the hierarchy (J'_α) , $\alpha \in Lim$, as follows:

$$Tr = \{(\alpha, \varphi(\mathbf{a})) : (J'_\alpha, \in, Tr \upharpoonright \alpha) \models \varphi(\mathbf{a}), \\ \varphi(\bar{x}) \in \mathcal{L}(\mathbf{aa}), \mathbf{a} \in J'_\alpha, \alpha \in Lim\},$$

where

$$Tr \upharpoonright \alpha = \{(\beta, \psi(\mathbf{a})) \in Tr : \beta \in \alpha \cap Lim\},$$

and

$$\begin{aligned} J'_0 &= \emptyset \\ J'_{\alpha+\omega} &= \text{rud}_{Tr}(J'_\alpha \cup \{J'_\alpha\}) \\ J'_{\omega\nu} &= \bigcup_{\alpha < \nu} J'_{\omega\alpha}, \text{ for } \nu \in Lim \\ C(\mathbf{aa}) &= \bigcup_{\alpha = \cup \alpha} J'_\alpha. \end{aligned}$$

Stationary tower forcing (see e.g. [8])

- Suppose δ is a Woodin cardinal².
- There is a forcing notion $Q = Q_{<\delta}$ such that $|Q| = \delta$ and such that if $G \subseteq Q$ is generic over V then in $V[G]$:
- δ is still a cardinal.
- There is an elementary embedding $j : V \rightarrow M$ where M is a transitive class such that $j(\omega_1) = \delta$. and such that $M^\omega \subseteq M$.

²A cardinal δ is **Woodin**, if for all $f : \delta \rightarrow \delta$ there is $\kappa < \delta$, closed under f , and $i : V \rightarrow M$ with critical point κ such that $V_{j(f(\kappa))} \subseteq M$.

C^* is forcing absolute, assuming PCW

- Suppose \mathbb{P} is a po-set.
- Let G be \mathbb{P} -generic.
- Choose a Woodin cardinal $\lambda > |\mathbb{P}|$.
- Let H_1 be generic for the stationary tower forcing $\mathbb{Q}_{<\lambda}$.
- In $V[H_1]$ there is a generic embedding $j_1 : V \rightarrow M_1$ such that $V[H_1] \models M_1^\omega \subseteq M_1$ and $j(\omega_1) = \lambda$.
- Hence $(C^*)^{V[H_1]} = (C^*)^{M_1}$ and³

$$j_1 : (C^*)^V \rightarrow (C^*)^{M_1} = (C^*)^{V[H_1]} = (C_{<\lambda}^*)^V.$$

- Now by elementarity $\text{Th}((C^*)^V) = \text{Th}((C_{<\lambda}^*)^V)$.

³ $C_{<\lambda}^*$ asks whether the cofinality of a linear order is $<\lambda$.

- Since $|\mathbb{P}| < \lambda$, λ is still Woodin in $V[G]$.
- Let H_2 be generic for $\mathbb{Q}_{<\lambda}$ over $V[G]$.
- Let $j_2 : V[G] \rightarrow M_2$ be the generic embedding.
- Now $V[G, H_2] \models M_2^\omega \subseteq M_2$ and $j_2(\omega_1) = \lambda$.
- Hence

$$j_2 : (C^*)^{V[G]} \rightarrow (C^*)^{M_2} = (C^*)^{V[G, H_2]} = (C^*_{<\lambda})^{V[G]} = (C^*_{<\lambda})^V,$$

- By elementarity $(C^*)^V \equiv (C^*_{<\lambda})^V \equiv (C^*)^{V[G]}$.

An important property of $C(\text{aa})$: Club Determinacy

- For all α :

$$(J'_\alpha, \in, Tr \upharpoonright \alpha) \models \forall \bar{x} [\text{aa} s \varphi(\bar{x}, \bar{t}, s) \vee \text{aa} s \neg \varphi(\bar{x}, \bar{t}, s)],$$

where $\varphi(\bar{x}, \bar{t}, s)$ is any formula in $L(\text{aa})$ and \bar{t} is a finite sequence of countable subsets of J'_α .

- CD **follows** from a proper class of Woodin cardinals [5].
- CD **follows** from PFA [5].

Theorem ([5])

- *Assuming Club Determinacy, every regular $\kappa \geq \aleph_1$ is measurable⁵ in $C(\text{aa})$.*
- *Suppose there are a proper class of Woodin cardinals. Then the first order theory of $C(\text{aa})$ is (set) forcing absolute.*

⁵We take, for a big α , all $X \subseteq \kappa$ in J'_α which in J'_α satisfy $\text{aa}(s \cap \kappa \in X)$

Proof of the forcing absoluteness of $C(\text{aa})$

- Suppose \mathbb{P} is a forcing notion and δ a Woodin cardinal $> |\mathbb{P}|$.
- Let $j : V \rightarrow M$ be the (generic) associated elementary embedding.
- Now $C(\text{aa}) \equiv (C(\text{aa}))^M = C(\text{aa}_\delta)$.
- Let $H \subseteq \mathbb{P}$ be generic over V and $j' : V[H] \rightarrow M'$.
- Again: $(C(\text{aa}))^{V[H]} \equiv (C(\text{aa}))^{M'} = (C(\text{aa}_\delta))^{V[H]}$.
- But $(C(\text{aa}_\delta))^{V[H]} = C(\text{aa}_\delta)$, since $|\mathbb{P}| < \delta$

Deeper into $C(aa)$

Deeper understanding of $C(aa)$ requires development of the theory of **aa-mice**.

Preparation for the definition of aa-premise

- We fix the following notation:

$$\tau_\xi = \{R_\infty, R_T, R_{T^*}\} \cup \{P_\eta : \eta < \xi\}, \quad \tau_\xi^- = \tau_\xi \setminus \{R_{T^*}\}.$$
- Here R_∞ and R_T are binary and R_{T^*}, P_η ($\eta < \xi$), are unary.
- We use $(P)_\xi$ to denote a sequence $\langle P_\eta : \eta < \xi \rangle$.

Example

The **canonical** example of an aa-premouse is

$$\mathcal{N} = (J'_\alpha, \in, Tr \upharpoonright \alpha, Tr_\alpha),$$

where $Tr_\alpha = \{\varphi(\mathbf{a}) : (\alpha, \varphi(\mathbf{a})) \in Tr\}$. Note that $\mathcal{N} \in C(\text{aa})$.

Definition

Let P^* be a **new** unary predicate symbol and $(P^*)^{M^*} = \{j(\mathbf{a}) : \mathbf{a} \in J_\alpha^T\}$. We let S^* consist of

$$\psi(P^*, [\varphi_1(s, x, \mathbf{a})], \dots, [\varphi_n(s, x, \mathbf{a})]),$$

where $\psi(s, x_1, \dots, x_n)$ is a τ -formula of $L(\text{aa})$, and

$$\text{aa}s\psi(s, f_{\varphi_1(s,x,\mathbf{a})}(s), \dots, f_{\varphi_n(s,x,\mathbf{a})}(s)) \in T^*.$$

Proposition

Let $\langle (M_\beta, E_\beta, T_\beta, T_\beta^*, (P^\beta)_\beta), j_{\beta,\gamma} : \beta \leq \gamma \leq \omega_1 \rangle$ be an aa-iteration of aa-mice. Then for all formulas $\varphi(\mathbf{a})$ of stationary logic in vocabulary $\tau_{\omega_1}^-$ and all $\mathbf{a} \in M_{\omega_1}$:

$$\varphi(\mathbf{a}) \in T_{\omega_1}^* \iff (M_{\omega_1}, E_{\omega_1}, T_{\omega_1}, (P^{\omega_1})_{\omega_1}) \models \varphi(\mathbf{a}).$$

Lemma

Suppose

$$(M_0, \in, T_0, T_0^*, (P)_0) \prec (J'_{\omega\alpha}, \in, Tr \upharpoonright \omega\alpha, Tr_{\omega\alpha}, (P')_0),$$

where α is a limit ordinal and M_0 is countable. Then M_{ω_1} **does not have new reals** over those in M_0 .

Theorem ([5])

If Club Determinacy holds, then **CH holds in $C(aa)$** .

Theorem ([5])

If club determinacy holds, there is a Δ^1_3 well-ordering of the reals in $C(aa)$. The reals form a countable Σ^1_3 -set.

Proof.

The canonical well-order \prec of $C(aa)$ satisfies:

$$x \prec y \iff \exists z \subseteq \omega \text{ (} z \text{ codes an aa-mouse } M \text{ such that } \\ x, y \in M \text{ and } M \models \text{ " } x \prec y \text{ ").}$$

The right hand side of the equivalence is Σ^1_3 and the claim follows. □

A summary of this lecture

- Stationary logic has a completeness theorem.
- Stationary logic gives rise to $C(\text{aa})$.
- Assuming large cardinals, $C(\text{aa})$ is forcing absolute, has measurable cardinals, and satisfies CH.

The next Lecture:

We learn:

- Why the **second order** constructible universe $C(\mathcal{L}^2)$ is the strongest?
- How to force the powerful **Henkin quantifier** constructible universe part ways with $C(\mathcal{L}^2)$?
- Why quantifying over **countable sets** gives a perfect (?) inner model?

Tutorial Lecture Three

We learn:

- Why the **second order** constructible universe $C(\mathcal{L}^2)$ is the strongest?
- How to force the powerful **Henkin quantifier** constructible universe part ways with $C(\mathcal{L}^2)$?
- Why quantifying over **countable sets** gives a perfect (?) inner model?

Gödel introduced HOD in 1946 [4].

Theorem ([11])

$C(\mathcal{L}^2) = \text{HOD}$.

$C(\mathcal{L}^2) \subseteq \text{HOD}$

- Let (L'_α) be the hierarchy building $C(\mathcal{L}^2)$.
- We use induction on α to show $L'_\alpha \in \text{HOD}$.
- Suppose $X \in L'_{\alpha+1}$, i.e. $X = \{a \in L'_\alpha : L'_\alpha \models \varphi(a, \vec{b})\}$, where $\varphi(x, \vec{y})$ is a second order formula in the language of set theory and $\vec{b} \in L'_\alpha$.
- Now $X = \{a : a \in L'_\alpha \wedge L'_\alpha \models \varphi(a, \vec{b})\}$, whence $X \in \text{HOD}$ as the truth predicate " \models " of second order logic is a set-theoretical predicate.

$C(\mathcal{L}^2) \supseteq \text{HOD}$

- Let $X \in \text{HOD}$.
- There is a first order $\varphi(x, \vec{y})$ and ordinals $\vec{\beta}$ such that for all a

$$a \in X \iff \varphi(a, \vec{\beta}).$$

- By Levy Reflection there is an α such that $X \subseteq V_\alpha$ and for all $a \in V_\alpha$

$$a \in X \iff V_\alpha \models \varphi(a, \vec{\beta}).$$

- We formalize the intuitive idea:
- Let $\theta(x, \vec{z})$ be a second order formula of the vocabulary $\{E, P\}$, E binary (“ \in ”), P unary (“the class of all ordinals”), such that:
- $(M, E, P) \models \theta(a^*, \vec{\beta}^*) \iff$ there is

$$\pi : (M, E, P) \cong (V_\delta, \in, \delta)$$

such that

$$(V_\delta, \in) \models \varphi(\pi(a^*), \pi(\vec{\beta}^*)).$$

Assume for a moment we have such a θ .

The following are equivalent, showing that $X \in L'_{\gamma+1}$:

- (1) $a \in X$
- (2) $L'_\gamma \models \exists M \exists E ((M, E, \alpha) \models \theta(a, \vec{\beta}) \wedge \text{TC}(\{a\}) \cup \alpha + 1 \cup \vec{\beta} \cup \{\vec{\beta}\} \subseteq M)$.

- (1) $a \in X$
- (2) $L'_\gamma \models \exists M \exists E ((M, E, \alpha) \models \theta(a, \vec{\beta}) \wedge \text{TC}(\{a\}) \cup \alpha + 1 \cup \vec{\beta} \cup \{\vec{\beta}\} \subseteq M)$.
- (1) \rightarrow (2) : Suppose $a \in X$. Thus $V_\alpha \models \varphi(a, \vec{\beta})$.
 - Let $M \subseteq L'_\gamma$ and $E \subseteq M^2$ s. t. $\alpha + 1, \text{TC}(a), \vec{\beta} \in M$ and there is an isomorphism $f : (V_\alpha, \in, \alpha, a, \vec{\beta}) \cong (M, E, P, a^*, \vec{\beta}^*)$. Remember $|L'_\gamma| \geq |V_\alpha|$.
 - We can assume $P = \alpha$, $a^* = a$ and $\vec{\beta}^* = \vec{\beta}$ by doing a partial Mostowski collapse⁶ for (M, E) , since $\alpha + 1, \text{TC}(a), \vec{\beta} \in M$.
 - Then $(M, E, \alpha) \models \varphi(a, \vec{\beta})$, whence $(M, E, \alpha) \models \theta(a, \vec{\beta})$, i.e. (2).

⁶In the Mostowski collapse everything that is transitive collapses onto itself.

- (1) $a \in X$
- (2) $L'_\gamma \models \exists M \exists E ((M, E, \alpha) \models \theta(a, \vec{\beta}) \wedge \text{TC}(\{a\}) \cup \alpha + 1 \cup \vec{\beta} \cup \{\vec{\beta}\} \subseteq M)$.
- (2) \rightarrow (1) : Suppose $M \subseteq L'_\gamma$, $P \subseteq M$, and $E \subseteq M^2$ such that $\text{TC}(\{a\}) \cup \alpha + 1 \cup \vec{\beta} \cup \{\vec{\beta}\} \subseteq M$ and $(M, E, P) \models \theta(a, \vec{\beta})$.
 - We may assume $E \upharpoonright \text{TC}(\{a\}) \cup \alpha + 1 \cup \vec{\beta} \cup \{\vec{\beta}\} = \in \upharpoonright \text{TC}(\{a\}) \cup \alpha + 1 \cup \vec{\beta} \cup \{\vec{\beta}\}$.
 - There is an isomorphism $\pi : (M, E, P) \cong (V_\alpha, \in, \alpha)$ such that $(V_\alpha, \in) \models \varphi(\pi(a), \pi(\vec{\beta}))$.
 - But $\pi(a) = a$ and $\pi(\vec{\beta}) = \vec{\beta}$.
 - So in the end $(V_\alpha, \in) \models \varphi(a, \vec{\beta})$. We have proved (1).

Construction of the formula θ

$\theta(x, z)$ is the conjunction, in the vocabulary $\{E, P\}$, of:

1. E is extensional and well-founded (i.e. there is no infinite sequence (a_n) such that $a_{n+1}Ea_n$ for all n .)
2. The class of the ordinals of the model is P .
3. If the cumulative hierarchy is built inside M using E as the ε -relation, then at successor stages M_{a+1} *every* subset of M_a is in M_{a+1} .
4. $\varphi(x, z)$ translated into the $\{E\}$ -vocabulary.

Note that for all β and $A \in \text{HOD}_1$:

- $\{\alpha < \beta : \text{cf}^V(\alpha) = \omega\} \in \text{HOD}_1$
- $\{(a, b) \in A^2 : |a|^V \leq |b|^V\} \in \text{HOD}_1$
- $\{\alpha < \beta : \alpha \text{ cardinal in } V\} \in \text{HOD}_1$
- $\{\alpha < \beta : (2^{|\alpha|})^V = (|\alpha|^+)^V\} \in \text{HOD}_1$

Lemma ([6])

1. $C^* \subseteq \text{HOD}_1$.
2. $C(Q_1^{MM, < \omega}) \subseteq \text{HOD}_1$.
3. *If 0^\sharp exists, then $0^\sharp \in \text{HOD}_1$ (by 1).*
 - Naturally, $\text{HOD}_1 = \text{HOD}$ is consistent, since we can assume $V = L$.
 - Also $\text{HOD}_1 \neq \text{HOD}$ is consistent relative to $\text{Con}(\text{ZF})$ ([6]).
 - Now we ask the same questions in the presence of large cardinals.

Theorem ([6, 9])

It is consistent, relative to the consistency of a supercompact cardinal that $\text{HOD} = \text{HOD}_1$ and there is also a supercompact cardinal.

We use [10] where it is proved that $V = \text{HOD}$ is consistent with a supercompact. The coding used in [10] gives also that $V = \text{HOD}_1$ is consistent with a supercompact.

Theorem ([6, 9])

It is consistent, relative to the consistency of a supercompact cardinal that for some λ :

$$\{\kappa < \lambda : \kappa \text{ weakly compact}\} \notin HOD_1,$$

and there is a supercompact cardinal.

Corollary

$HOD \neq HOD_1$ is consistent with a supercompact.

- [10] shows that one can start with any model with a supercompact cardinal and force $V = \text{HOD}$ preserving the supercompact cardinal.
- It is enough to code subsets of cardinals κ such that $\kappa = \beth_{\kappa}$.
- Menas codes subsets into the relation $2^{\aleph_{\alpha+1}} = \aleph_{\alpha+\omega+3}$.
- A set $X \subseteq \aleph_{\nu}$ ends up being definable as

$$\{\alpha < \aleph_{\nu} : 2^{\aleph_{e(\nu)+\alpha+1}} = \aleph_{e(\nu)+\alpha+3}\},$$

where $e(\nu)$ is the ν th ordinal α such that $\alpha = \beth_{\alpha}$.

- Hence his model satisfies $V = \text{HOD}_1$.

To be done on the blackboard

- As in [7], we define in $V^{\mathbb{P}_n}$ a forcing \mathbb{S}_n to be the canonical forcing which introduces a κ_n homogeneous Soulin tree. It kills the weak compactness of κ_n .
- Let \mathbb{T}_n be the forcing which introduces a branch through the tree forced by \mathbb{S}_n .
- $\mathbb{S}_n * \mathbb{T}_n$ is forcing equivalent to \mathbb{D}_{κ_n} .
- If we force with \mathbb{T}_n over $V^{\mathbb{P}_n * \mathbb{S}_n}$, we regain the weak compactness of κ_n .
- Also a generic object for \mathbb{D}_{κ_n} introduces a generic object for \mathbb{T}_n .

To be done on the blackboard

- Let $V_1^* = V_0^{\mathbb{Q}}$.
- Let G_n be the generic filter in \mathbb{P}_n , introduced by \mathbb{Q} .
- The model V_3^* is the model one gets from V_1^* by forcing over it with the full support product of \mathbb{D}_{κ_n} . (\mathbb{D}_{κ_n} is as realized according to G_n .)
- Let $H_n \subseteq \mathbb{D}_{\kappa_n}$ be the generic filter introduced by this forcing.
- Note that V_3^* can also be obtained from V_0 by forcing with \mathbb{Q} . In particular both in V_1 and in V_3^* the cardinals κ_n are weakly compact for every $n < \omega$. Let V_3 be an extension of V_3^* by adding a Cohen real $a \subseteq \omega$.
- Let \mathbb{A} be the Cohen forcing on ω .
- Let $V_1 = V_1^*(a)$. Both V_1 and V_3 are obtained by forcing over V_0 with $\mathbb{Q} \times \mathbb{A}$ which is a homogenous forcing notion. Hence $HOD^{V_1} = HOD^{V_3}$.
- Again we did not kill the weak compactness of the cardinals κ_n .

To be done on the blackboard

- By the standard arguments analysing the power-set of a cardinal δ under a forcing which is the product of a forcing of size $\mu < \delta$, a forcing of size δ which is μ^+ -distributive, and a forcing which is δ^+ -distributive:
- For $n < \omega$ $P(\kappa_n)^{V_2} = P(\kappa_n)^{W_n}$.
- $V_2 \models a = \{n < \omega \mid \kappa_n \text{ is weakly compact}\}$.
- $a \in HOD^{V_2}$.

To be done on the blackboard

Lemma

The following are equivalent

1. $(M \models \Phi(\vec{b}))^{V_1}$.
2. $(M \models \Phi(\vec{b}))^{V_3}$.
3. $(M \models \Phi(\vec{b}))^{V_2}$.

To be done on the blackboard

- Without loss of generality, $\Phi(\vec{x})$ has the form $\exists X \Psi(X, \vec{x})$.
- Both V_1 and V_3 are obtained from V_0 by forcing over V_0 with $\mathbb{Q} \times \mathbb{A}$. This forcing is homogeneous.
- M and all the elements of the vector \vec{b} are in V_0 .
- So (1) is clearly equivalent to (2).
- Suppose that $(M \models \Phi(\vec{b}))^{V_2}$. Let $Z \subseteq M$ be the witness for the existential quantifier of Φ . Then $(M \models \Psi(Z, \vec{b}))^{V_2}$. But all the quantifiers of Ψ are first order, so $(M \models \Psi(Z, \vec{b}))^{V_3}$.
- So (3) implies (2), and hence (1).
- For the other direction, if $(M \models \Phi(\vec{b}))^{V_3}$, then we know that $(M \models \Phi(\vec{b}))^{V_1}$. Let $Z \in V_1$ satisfy $M \models \Psi(Z, \vec{b})$. So $(M \models \Psi(Z, \vec{b}))^{V_2}$, and therefore $(M \models \Phi(\vec{b}))^{V_2}$.

To be done on the blackboard

- It follows from the lemma that every Σ_1^1 formula defines the same subset of M in V_1 , V_2 and V_3 .
- It follows that $L_{\alpha+1}^1 = L_{\alpha+1}^2 = L_{\alpha+1}^3$.
- This proves the lemma and the theorem.
- **Conclusion:** Large cardinals cannot decide the question “ $HOD = HOD_1$?”

Quantifying over countable sets⁹: $C^2(\omega)$

- $ZFC \vdash C^2(\omega) \subseteq HOD^{C_{\omega_1 \omega_1}}$. Hence $Th(C^2(\omega))$ is **forcing absolute**, assuming a proper class of Woodin limits of Woodins. (Woodin)
- $V = C^2(\omega)$ implies there are no measurable cardinals.
- ω_1^V is **strongly Mahlo** in $C^2(\omega)$, assuming a Woodin limit of Woodins. [9]
- $C^2(\omega)$ contains, for every n , the inner model with n **Woodin cardinals**.
- If there is a proper class of Woodin limits of Woodin cardinals, then every regular cardinal is **measurable** in $C^2(\omega, aa)$ and $C^2(\omega, aa)$ is forcing absolute [9].
- CH?

⁹Joint work with Menachem Magidor.

Recent developments

1. Assume a proper class of Woodin cardinals. The reals of $C(aa)$ (and also of C^*) are in M_1 . $C(aa)$ has no inner model with a Woodin cardinal. (Magidor-Schindler)
2. Assume Club Determinacy. Then **Ultrapower Axiom** and **GCH** hold in $C(aa)$. (Goldberg-Steel)
3. Assume Club Determinacy. If κ is regular in $C(aa)$ with $cof(\kappa) \geq \omega_2^V$, then $o(\kappa)^{C(aa)} \geq 2$. Moreover, then $o(\omega_3^V)^{C(aa)} \geq 3$. (Goldberg-Rajala)
4. If $V = L^\mu$, then $V = C(aa)$. (Goldberg, Magidor, Schindler, Steel)
5. There is ongoing investigation on what kind of **mice** can be found inside models such as C^* and $C(aa)$. (Goldberg, Magidor, Schindler, Steel)

A summary of this lecture

- The second logic version of L is actually HOD.
- The Henkin quantifier version of L , called HOD₁, can be smaller than HOD
- Quantifying over countable subsets yields an inner model $C^2(\omega, aa)$ which is, assuming large cardinals, forcing absolute and has measurable cardinals. Does it have CH ?

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Thank you!